# Enriched duality in double categories

Christina Vasilakopoulou

National Technical University of Athens, Greece

Workshop on Categorical Logic and Higher Categories

University of Manchester

#### Outline

- Background
- 2. Enriched duality in monoidal double categories
- Oplax monoidal double structure
- Further directions

The category Alg of k-algebras is enriched in Coalg of k-coalgebras.

#### Monoids and comonoids

Suppose  $(\mathcal{V}, \otimes, I)$  is monoidal category.

▶ A monoid is an object A together with maps  $\mu$ :  $A \otimes A \to A$  and  $\eta$ :  $I \to A$  which are associative and unital

- ▶ Dually, a *comonoid* is an object C together with maps  $\delta \colon C \to C \otimes C$  and  $\epsilon \colon C \to I$  which are coassociative and counital.
- $\star$  These form categories Mon and Comon, with maps preserving structure.

When  $(\mathcal{V}, \otimes, I, \sigma)$  is symmetric, Mon and Comon are monoidal with

$$C \otimes D \xrightarrow{\delta \otimes \delta} C \otimes C \otimes D \otimes D \xrightarrow{1 \otimes \sigma \otimes 1} C \otimes D \otimes C \otimes D$$

■ If  $\mathcal V$  is monoidal closed, induced [-,-]: Comon $(\mathcal V)^{\mathrm{op}} \times \mathrm{Mon}(\mathcal V) \to \mathrm{Mon}(\mathcal V)$  makes [C,A] into a monoid via *convolution* 

$$[C,A] \otimes [C,A] \otimes C \xrightarrow{1 \otimes \delta} [C,A] \otimes [C,A] \otimes C \otimes C \xrightarrow{1 \otimes \sigma \otimes 1} [C,A] \otimes C \otimes [C,A] \otimes C \xrightarrow{\downarrow \text{ev} \otimes \text{ev}} A \otimes A \xrightarrow{\downarrow \mu} A \otimes A$$

\* In  $\operatorname{Vect}_k$ , linear dual  $C^* = \operatorname{Hom}_k(C, k)$  for a k-coalgebra is a k-algebra – while  $A^*$  for a k-algebra is a k-coalgebra only if it is finite dimensional.

Sweedler dual 'fixes' that:  $A^o = \{ f \in A^* \mid kerf \text{ contains cofinite ideal} \}$  is a k-coalgebra, and  $Alg(A, C^*) \cong Coalg(C, A^o)$ .

More generally, there exists universal measuring k-coalgebra with  $Alg(A, Hom_k(C, B)) \cong Coalg(C, P(A, B))$  – so  $A^o = P(A, k)$ .

Moving to general context of symmetric monoidal closed categories, local presentability gives passage from  $\text{Vect}_k$  to  $(d)g\text{Vect}_k$ ,  $\text{Mod}_R$  etc.

Suppose  $\mathcal V$  is a symmetric monoidal closed and locally presentable category. There is a parameterized adjunction between

[-,-]: Comon<sup>op</sup> 
$$\times$$
 Mon  $\rightarrow$  Mon convolution  $P(-,-)$ : Mon<sup>op</sup>  $\times$  Mon  $\rightarrow$  Comon universal measuring

- In Set, P(A, B) is Mon(A, B); in  $Vect_k$ , it contains k-algebra maps as grouplike elements; in  $dgVect_k$ , it relates to bar-cobar adjunction.
- \* Convolution [-,-] is an action of the monoidal Comon on Mon.

An adjoint of an action  $\bullet \colon \mathcal{V} \times \mathcal{C} \to \mathcal{C}$  gives rise to a  $\mathcal{V}$ -enriched structure on  $\mathcal{C}$ .

The category Mon is enriched in the monoidal Comon.

### From monoidal to double categories

One to many objects: generalize from monoids in V, to V-categories.

What about comonoids? Opcategories= $\mathcal{V}^{\mathrm{op}}$ -categories not as convenient, formally. . . identify common framework!

ightharpoonup A double category  $\mathbb{D}_0$  has object category  $\mathbb{D}_0$  (0-cells & vertical 1-cells), arrow category  $\mathbb{D}_1$  (horizontal 1-cells & 2-maps)

$$\begin{array}{ccc}
X & \xrightarrow{A} & Y \\
f \downarrow & \psi \alpha & \downarrow g \\
Z & \xrightarrow{B} & W
\end{array}$$

and  $\mathbb{D}_0 \xrightarrow{\mathbf{1}} \mathbb{D}_1$ ,  $\mathbb{D}_1 \overset{\mathfrak{s}}{\underset{t}{\Longrightarrow}} \mathbb{D}_0$ ,  $\mathbb{D}_1 \times_{\mathbb{D}_0} \mathbb{D}_1 \overset{\circ}{\to} \mathbb{D}_1$  + coherent isos.

0-cells, horizontal 1-cells, globular 2-maps make horizontal bicategory  $\mathcal{H}(\mathbb{D})$ .

ightharpoonup A monad in a double category  $\mathbb D$  is  $A\colon X \dashrightarrow X$  with associative, unital

Dually, comonad  $C: X \longrightarrow X$ . These form categories  $Mnd(\mathbb{D})$  &  $Cmd(\mathbb{D})$ .

\* Morphisms are different than those for (co)monads in bicategories.

For 
$$\mathbb{D} = \mathcal{V}$$
-Mat of sets, functions and  $\mathcal{V}$ -matrices  $S \colon X \to Y$  i.e.  $\{S(x,y)\} \in \mathcal{V}$  with  $(S \circ T)(x,z) = \sum_y T(x,y) \otimes S(y,z)$ ,  $\mathsf{Mnd}(\mathcal{V}\text{-Mat}) = \mathcal{V}\text{-Cat}$  and  $\mathsf{Cmd}(\mathcal{V}\text{-Mat}) = \mathcal{V}\text{-Cocat}$ .

A  $\mathcal{V}$ -cocategory comes with cocomposition  $C(x,z) \to \sum_y C(x,y) \otimes C(y,z)$  and coidentities  $C(x,x) \to I$ , coassociative and counital.

For  $\mathbb D$  a  $(\dots)$  double category,  $\mathsf{Mnd}(\mathbb D)$  is enriched in  $\mathsf{Cmd}(\mathbb D)$ .

- Fibrant: vertical 1-cells f turn to horizontal, companion  $\hat{f}$  & conjoint  $\check{f}$ . In  $\mathcal{V}$ -Mat,  $f: X \to Y$  gives matrices  $\hat{f}(x,y) = \check{f}(y,x) = \begin{cases} I \text{ if } fx = y \\ 0 \text{ if } fx \neq y \end{cases}$ 
  - $\mathsf{Mnd}(\mathbb{D}) \to \mathbb{D}_0$  is a fibration; reindexing  $\check{f} \circ \circ \hat{f} \colon \mathsf{Mnd}(\mathbb{D})_Y \to \mathsf{Mnd}(\mathbb{D})_X$ . Dually,  $\mathsf{Cmd}(\mathbb{D}) \to \mathbb{D}_0$  is an optibration.
- Monoidal:  $\mathbb{D}_0$  &  $\mathbb{D}_1$  monoidal,  $(N \circ M) \otimes (N' \circ M') \cong (N \otimes N') \circ (M \otimes M')$ . In  $\mathcal{V}$ -Mat,  $(X \otimes Y) = X \times Y$  &  $(S \otimes T)((x, y), (z, w)) = S(x, z) \otimes T(y, w)$ .

 $\mathsf{Mnd}(\mathbb{D})$  and  $\mathsf{Cmd}(\mathbb{D})$  are monoidal, with  $C \otimes D$  comonad via  $C \otimes D \to (C \circ C) \otimes (D \circ D) \cong (C \otimes D) \circ (C \otimes D)$ .

■ Monoidal closed : lax double functor  $H \colon \mathbb{D}^{\mathrm{op}} \times \mathbb{D} \to \mathbb{D}$  such that  $\mathbb{D}_0$  &  $\mathbb{D}_1$  monoidal closed,  $\mathfrak{s}, \mathfrak{t}$  maps of adjunctions.

In 
$$\mathcal{V}$$
-Mat,  $[X, Y] = Y^X$  and  $H(S, T)(f, g) = \prod_{x,y} [S(x, y), T(fx, gy)].$ 

H induces functor  $\mathsf{Cmd}(\mathbb{D})^{\mathrm{op}} \times \mathsf{Mnd}(\mathbb{D}) \to \mathsf{Mnd}(\mathbb{D})$ , an action. *convolution* 

■ Locally presentable :  $\mathbb{D}_0$  &  $\mathbb{D}_1$  locally presentable,  $\mathfrak{s},\mathfrak{t}$  cocontinuous right adjoints, -  $\circ$  - accessible in each variable.

In 
$$\mathcal{V}$$
-Mat, Set is I. p. &  $\mathcal{V}$ -Mat<sub>1</sub> is too as pullback 
$$\begin{array}{c} \mathcal{V}\text{-Mat}_1 \to \mathsf{Fam}(\mathcal{V})\\ (\mathfrak{s},\mathfrak{t})\downarrow & \downarrow\\ \mathsf{Set}^2 \stackrel{\times}{\longrightarrow} \mathsf{Set} \end{array}$$

by the Limit Theorem.

Induced H has adjoint  $\mathsf{Mnd}(\mathbb{D})^{\mathrm{op}} \times \mathsf{Mnd}(\mathbb{D}) \to \mathsf{Cmd}(\mathbb{D})$ . univ. measuring

 $\diamond$  Obtain enrichment of Mnd( $\mathbb{D}$ ) in Cmd( $\mathbb{D}$ )!

## Moving to other contexts

- $\star$  For  $\mathcal V$  symmetric monoidal closed & locally presentable,  $\mathbb D=\mathcal V\text{-}\mathbb M$ at is a fibrant, monoidal closed & locally presentable double category.
  - Enrichment of  $\mathcal{V}$ -categories in  $\mathcal{V}$ -cocategories

Goal: employ/extend theory to apply to  $\mathbb{D} = \mathcal{V}\text{-}\mathbb{S}\text{ym}$ 



- · objects are sets  $X, Y, \ldots$
- · vertical 1-cells are functions  $f, g, \ldots$
- · horizontal 1-cells are coloured symmetric sequences  $M: SX^{\mathrm{op}} \times Y {
  ightharpoonup} \mathcal{V}$
- · 2-maps are  $M(x_1, \ldots, x_n; y) \rightarrow N(fx_1, \ldots, fx_n; gy)$

Horizontal composition is generalization of substitution for species.

■ V-Sym is fibrant; but not monoidal double anymore!

### Oplax monoidal double structure

► A double category is *oplax* monoidal with comparison maps

$$(N \circ M) \otimes (N' \circ M') \rightarrow (N \otimes N') \circ (M \otimes M'), \quad 1_X \otimes 1_{X'} \rightarrow 1_{X \otimes X'}$$
  
 $I_1 \rightarrow I_1 \circ I_1, \quad I_1 \rightarrow 1_I$ 

satisfying coherence axioms.

An oplax monoidal double category with a single object and vertical arrow is precisely a duoidal category.

End result:  $\mathcal{V}$ -Sym is an oplax monoidal double category. How? As a Kleisli-type structure on  $\mathcal{V}$ -Prof with induced monoidality.

ightharpoonup A (vertical) double monad is a double functor  $T: \mathbb{D} \to \mathbb{D}$  with transformations  $m: TT \Rightarrow T, e: 1 \Rightarrow T$  with

- ▶ It is special when  $\hat{m}_X$ ,  $\hat{e}_X$  exist, and transposes of  $m_M$ ,  $e_M$  are invertible.
- lacksquare Each special double monad  $T\colon \mathbb{D} o\mathbb{D}$  gives double category  $\mathbb{K}\mathsf{I}(T)$ 
  - $\cdot \mathbb{K}I(T)_0$  is  $\mathbb{D}_0$
  - $\cdot$   $M: X \rightsquigarrow Y$  are horizontal  $M: X \longrightarrow TY$  in  $\mathbb{D}$

 $\cdot \text{ horizontal composition is } X \stackrel{M}{\longrightarrow} TY \stackrel{TN}{\longrightarrow} TTZ \stackrel{\hat{m}_Z}{\longrightarrow} TZ$ 

ightharpoonup A double monad T on a monoidal double category  $\mathbb D$  is pseudomonoidal when T lax monoidal and m, e pseudomonoidal

For T a pseudomonoidal special double monad, if lax structure maps  $I \to TI$ ,  $TX \otimes TY \xrightarrow{\tau} T(X \otimes Y)$  have companions,  $\mathbb{K}I(T)$  is an oplax monoidal double category.

\* Induced tensor is  $M \boxtimes N = X \otimes Z \xrightarrow{M \otimes N} TY \otimes TW \xrightarrow{\hat{\tau}} T(Y \otimes W)$ .

### Coloured symmetric sequences

 $\star$  'Free symmetric strict monoidal category 2-monad'  $S \colon \mathcal{V}\text{-Cat} \to \mathcal{V}\text{-Cat}$ 

with 
$$S_n(C)((x_1,\ldots,x_n),(y_1,\ldots,y_n)) = \sum_{\sigma} \prod_{1\leq i\leq n} C(x_{\sigma(i)},y_i)$$

extends to such a double monad on the monoidal double category  $\mathcal{V}$ - $\mathbb{P}$ rof (for  $\mathcal{V}$  cartesian monoidal).

Kleisli double category is  $\mathcal{V}$ - $\mathbb{C}$ at $\mathbb{S}$ ym of categorical symmetric sequences  $M: SX^{\mathrm{op}} \times Y \rightarrow \mathcal{V}$ , with 'discrete' case  $\mathcal{V}$ -Sym.

Oplax monoidal structure is many-object arithmetic product of species

$$(M\boxtimes N)(\vec{a},(x,z))=\int^{\vec{y},\vec{w}}S(Y\times W)(\vec{a},\vec{y}\boxtimes\vec{w})\times M(\vec{y},x)\times N(\vec{w},z)$$

#### Some future directions

- ► Extend previous results from monoidal double to <u>oplax</u> monoidal double categories : do we still obtain enrichment of monads in comonads?
- ightharpoonup Further explore  $\mathcal{V} ext{-}\mathbb{S}$ ym: is it monoidal closed and locally presentable as a double category?

[One-object case] If  ${\mathcal V}$  is symmetric monoidal closed and loc presentable,

- positive operads are enriched in positive cooperads, if  ${\cal V}$  has biproducts;
- ullet symmetric operads are enriched in symmetric cooperads, if  ${\cal V}$  is cartesian.

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\mathsf{Mod}(\mathbb{D}) \overset{\mathrm{enriched}}{\longrightarrow} \mathsf{Comod}(\mathbb{D})
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#### Thank you for your attention!



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